

# Exploiting Tropical Algebra in Numerical Linear Algebra

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**Householder Symposium XIX, June 2014.**

# Tropical Semirings

By “**tropical**” we refer to a semiring in which the addition operation is **min** or **max**.

In this talk, consider **max-plus semiring**  $(\mathbb{R}_{\max}, \oplus, \otimes)$ , where  $\mathbb{R}_{\max} = \mathbb{R} \cup \{-\infty\}$ ,

$$a \oplus b = \max(a, b), \quad a \otimes b = a + b, \quad \forall a, b \in \mathbb{R}_{\max},$$

and additive and multiplicative identities  $-\infty$  and  $0$ :

$$a \oplus -\infty = a, \quad a \otimes 0 = a.$$

**Tropical algebra** is the tropical analogue of linear algebra, working with matrices with entries in  $\mathbb{R}_{\max}$ . If  $A, B \in \mathbb{R}_{\max}^{n \times n}$ ,

$$(A \oplus B)_{ij} = a_{ij} \oplus b_{ij}, \quad (A \otimes B)_{ij} = \bigoplus_{k=1}^n a_{ik} \otimes b_{kj}.$$

# When/How Can Tropical Algebra Help NLA?

**When?** Tropical algebra can help NLA when there are **large variations** in the magnitude of the data.

**How?** By providing **order of magnitude approximation** to roots, modulus of ei'vals and singular values.

- Offer good starting points for iterative algorithms.
- Can help to reduce condition numbers/backward errors.

error in solution  $\lesssim$  condition number  $\times$  backward error.

# "Tropicalization" of Linear Algebra Problems

- ▶ We use **valuations** (provide a measure of size or multiplicity of elements of the field).

- E.g.,  $x \in \mathbb{C} \mapsto \mathcal{V}(x) = \log |x| \in \mathbb{R}_{\max}$  ( $\log 0 = -\infty$ ).

- When  $|a| \gg |b|$  or  $|a| \ll |b|$  with  $a, b \in \mathbb{C}$ ,

$$\begin{aligned} \mathcal{V}(a + b) &= \log |a + b| & \mathcal{V}(ab) &= \log |ab| \\ &\approx \max(\log |a|, \log |b|) & &= \log |a| + \log |b| \\ &= \mathcal{V}(a) \oplus \mathcal{V}(b), & &= \mathcal{V}(a) \otimes \mathcal{V}(b). \end{aligned}$$

- ▶ Tropicalized linear algebra problems
  - can be easier/cheaper to solve and,
  - does not suffer much from numerical instabilities.

# Scalar Polynomials: Classical/Max-Plus

“**Tropicalize**”  $p(x) = \sum_{i=0}^d a_i x^i$ ,  $a_i \in \mathbb{C}$ , i.e., construct

$$\text{tp}(x) = \bigoplus_{i=0}^d \log |a_i| \otimes x^{\otimes i} = \max_{0 \leq i \leq d} (\log |a_i| + ix).$$

Let  $\alpha_1 < \dots < \alpha_p$  be roots of  $\text{tp}$  with  $\alpha_j$  of multiplicity  $m_j$ .

## Theorem (Sharify'11)

*If  $\max(\alpha_j - \alpha_{j-1}, \alpha_{j+1} - \alpha_j) \geq \log 9 \approx 2.2$  for  $1 \leq j \leq p$  then  $p(x)$  has exactly  $m_j$  roots in the annulus*

$$A(x) = \{x \in \mathbb{C} : \frac{1}{3} \exp(\alpha_j) \leq |x| \leq 3 \exp(\alpha_j)\}.$$

**Max-plus roots of  $\text{tp}(x)$  offer order of magnitude approx. to roots of  $p$  as long as the  $\alpha_j$  are well separated.**

# Extension to Matrix Polynomials

Let  $P(\lambda) = \sum_{i=0}^d A_i \lambda^i \in \mathbb{C}[\lambda]^{n \times n}$  and  $tp(x) = \bigoplus_{i=0}^d \log \|A_i\| \otimes x^{\otimes i}$

with max-plus roots  $\alpha_1 < \dots < \alpha_q$ ,  $\alpha_j$  of multiplicity  $m_j$ .

$k_0 < \dots < k_q$ : corresponding indices in Newton polygon.

## Theorem (Noferini, Sharify, T' 14)

*If  $\alpha_\ell - \alpha_{\ell-1} \geq 2 \log(1 + 2\kappa(A_{k_\ell}))$ ,  $\ell = j - 1, j$  then  $P(\lambda)$  has exactly  $nm_j$  ei'vals inside the annulus*

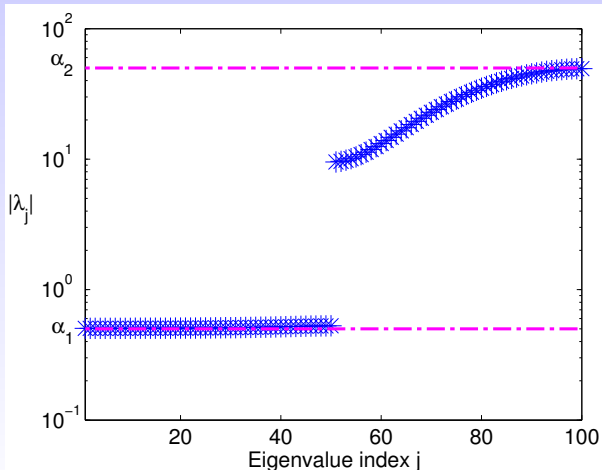
*$\mathcal{A}((1 + 2\kappa(A_{k_{j-1}}))^{-1} \exp(\alpha_j), (1 + 2\kappa(A_{k_j})) \exp(\alpha_j))$ .*

**For  $A_{k_{j-1}}, A_{k_j}$  well conditioned and  $\alpha_{j-1}, \alpha_j, \alpha_{j+1}$  sufficiently well separated,  $P$  has  $nm_j$  ei'vals of modulus close to  $\exp(\alpha_j)$ .**

# Illustration with Spring Problem

spring: quadratic matrix polynomial from NLEVP.

$$\kappa_2(\mathbf{A}_j) \leq 5, j = 0: 2.$$



# Use of Max-Plus Roots in NLA

- ▶ Max-plus roots used to select **starting points** in the Ehrlich-Aberth method for polynomial eigenproblems. [Bini, Noferini, Sharify'13].
- ▶ Useful for Betcke's **diagonal scaling** aimed at improving the conditioning of P's ei'vals near  $\omega$ .
- ▶ Define **eigenvalue parameter scalings** ( $\lambda = \alpha_j \mu$ ) for polynomial eigenvalue solvers based on linearizations.

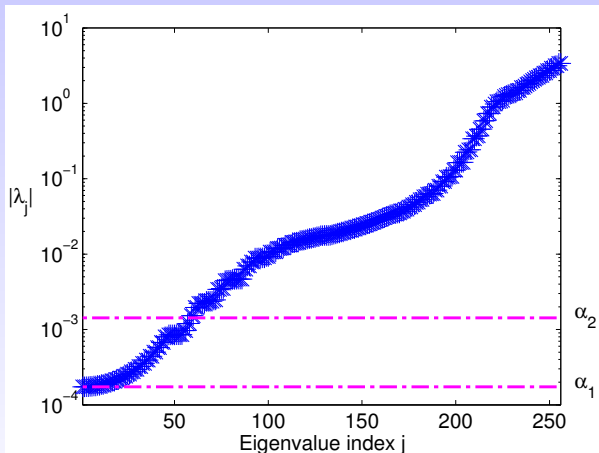
$$\tilde{P}(\mu) := \delta^{-1} P(\alpha_j \mu), \quad \delta = \|A_{k_{j-1}}\| \exp(k_{j-1} \alpha_j).$$

- Allow computation of ei'pairs with small b'err for  $|\lambda|$  near  $\alpha_j$ .
- Linearization process does not affect ei'val condition number of ei'vals near  $\alpha_j$ .



# Orr-Sommerfeld Problem

Quartic matrix polynomial from NLEVP collection.



Here e'vals of large magnitude are not captured by the max-plus roots.

# Max-Plus Eigenvalues

- The **max-plus ei'vals** of  $M \in \mathbb{R}_{\max}^{n \times n}$  are the roots of

$$\chi_M(\lambda) = \text{perm}(M \oplus \lambda \otimes I).$$

Here  $I$  is the identity matrix in  $\mathbb{R}_{\max}^{n \times n}$ .

- The  $n$  max-plus ei'vals of  $M$  can be computed in  $O(zn)$  ops, where  $z = \text{nnz}(M)$  using a network flow algorithm [Gassner & Klinz'10]
- The max-plus ei'vals of a max-plus matrix polynomial

$$P(\lambda) = P_0 \oplus P_1 \otimes \lambda \oplus \cdots \oplus P_d \otimes \lambda^{\otimes d}, \quad P_j \in \mathbb{R}_{\max}^{n \times n}$$

are the max-plus roots of  $\chi_P(\lambda) = \text{perm}(P(\lambda))$ .

# Matrices Depending on a Parameter

Let  $A(t)_{ij} = b_{ij} \exp(m_{ij}t)$ ,  $B \in \mathbb{C}^{n \times n}$ ,  $M \in \mathbb{R}_{\max}^{n \times n}$  ( $\exp(-\infty) = 0$ ).

Use valuation  $\mathcal{V}[f(t)] = \lim_{t \rightarrow \infty} \frac{\log |f(t)|}{t} \implies \mathcal{V}[A] = M$ .

## Theorem (Akian, Bapat, Gaubert' 04)

For all  $M \in \mathbb{R}_{\max}^{n \times n}$  and all **generic**  $B \in \mathbb{C}^{n \times n}$ , the ei'vals  $\lambda_1(t), \dots, \lambda_n(t)$  of nonsing.  $A(t)$  satisfy

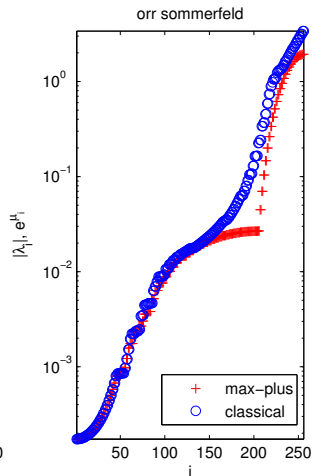
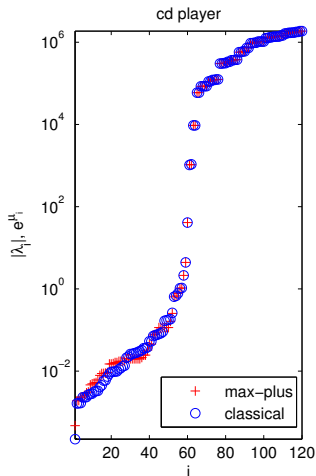
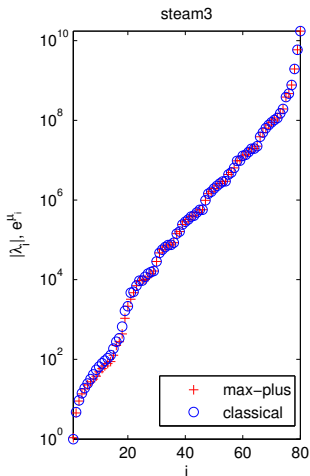
$$\mathcal{V}[\lambda_i(t)] = \lim_{t \rightarrow \infty} \frac{\log |\lambda_i(t)|}{t} = \mu_i,$$

where  $\mu_1, \dots, \mu_n$  are max-plus ei'vals of  $M$ 's.

For  $A \in \mathbb{C}^{n \times n}$  use  $\mathcal{V}[x] = \log |x| \implies (\mathcal{V}[A])_{ij} = \log |a_{ij}|$ .

# Classical/Max-Plus Eigenvalues

steam 3 from Florida sparse,  
cd\_player, orr\_sommerfeld from the NLEVP.



# Max-Plus Singular Values

Let  $M \in \mathbb{R}_{\max}^{n \times n}$ ,  $B \in \mathbb{C}^{n \times n}$ ,  $A(t) = (a_{ij}(t))$ ,  $a_{ij}(t) = b_{ij} \exp(m_{ij}t)$ .

## Theorem (De Schutter, De Moor' 02)

Let  $A(t) = U(t)\Sigma(t)V(t)$  be the analytic SVD of  $A(t)$  with  $\Sigma = \text{diag}(\sigma_1(t), \dots, \sigma_n(t))$ ,  $\sigma_j(t) = 0$ ,  $j = n - k + 1 : n$ . Then for all  $G$  and all **generic**  $B$

$$\lim_{t \rightarrow \infty} \frac{\log \sigma_i(t)}{t} =: s_i, \quad i = 1 : n - k$$

exists and is independent of the choice of  $B$ .

## Definition

The **max-plus singular values** of  $M$  are  $s_1, \dots, s_n$ , with  $s_1, \dots, s_{n-k}$  defined as above and  $s_{n-k+1}, \dots, s_n = -\infty$ .

# Max-Plus Singular Values (cont.)

## Theorem (Hook' 14)

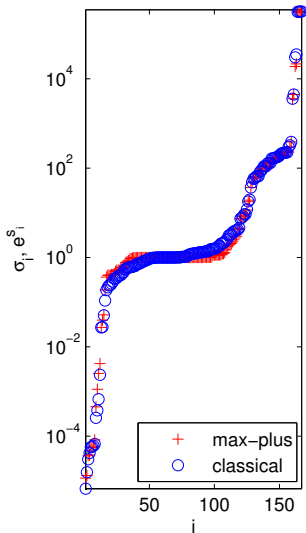
The **max-plus singular values** of  $M \in \mathbb{R}_{\max}^{n \times n}$  are the max-plus ei'vals of the max-plus pencil,

$$M \oplus \sigma \otimes 0.$$

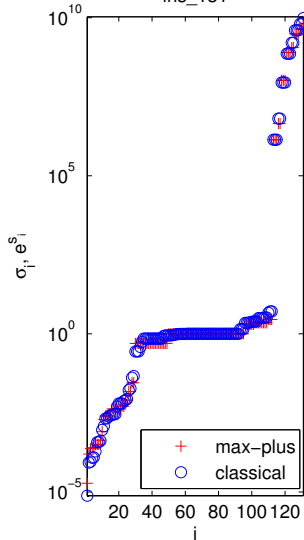
- Characterization extends to rectangular case.
- Max-plus singular values and e'vals of symmetric max-plus matrices are equal.
- If  $A \in \mathbb{C}^{n \times n}$  with  $i$ th singular value  $\sigma_i$  we expect  $\exp(s_i) \approx \sigma_i$ , where  $s_i$  is the  $i$ th max-plus singular value of  $M = (\log |a_{ij}|) \in \mathbb{R}_{\max}^{n \times n}$ .

# Matrices from Florida Sparse Collection

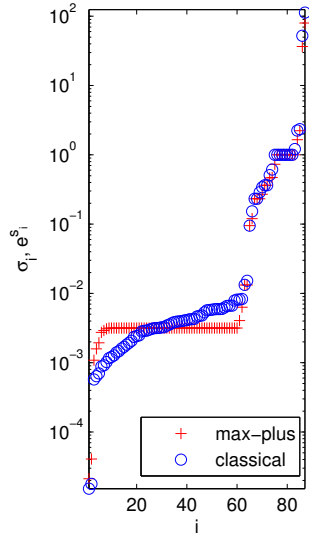
west0167



Ins\_131



d\_dyn



# Hungarian Pair

**Optimal assignment problem** for  $M \in \mathbb{R}_{\max}^{n \times n}$ : compute

$$\text{perm}(M) = \max_{\pi \in P(n)} \sum_{i=1}^n m_{\pi(i), i}.$$

Can be expressed as a **linear programming problem** (LPP)

$$\max \left\{ \sum_{i,j=1}^n m_{ij} d_{ij} : d_{ij} > 0, \sum_{j=1}^n d_{ij} = \sum_{j=1}^n d_{ji} = 1 \forall i \right\},$$

with **dual problem**

$$\min \left\{ \sum_{i=1}^n u_i + v_i : u, v \in \mathbb{R}^n : m_{ij} - u_i - v_j \leq 0 \right\}.$$

An **Hungarian pair** is a solution  $(u, v)$  to the dual LPP.



# Hungarian Scaling

For  $A \in \mathbb{C}^{n \times n}$  let  $M = \mathcal{V}[A] \in \mathbb{R}_{\max}^{n \times n}$  with  $\mathcal{V}[x] = \log |x|$ .

Let  $(u, v)$  be a Hungarian pair for  $M$ .

Define diagonal matrices  $D_u, D_v \in \mathbb{R}^{n \times n}$  by

$$(D_u)_{ii} = \exp(-u_i), \quad (D_v)_{ii} = \exp(-v_i).$$

The **Hungarian scaling** of  $A$  is  $D_u A D_v =: H$  ( $|h_{ij}| \leq 1$ ).

- ▶ Implemented in HSL-MC64 (together with some reordering).
- ▶ Commonly used before solving highly indefinite and nonsymmetric linear systems.
- ▶ Can show  $\kappa_2(H) \leq n \min_{D_1, D_2 \in \mathcal{D}_n} \kappa_2(D_1 A D_2)$  for normal  $A$  (tends to also be the case for nonnormal  $A$ ).

# Numerical Experiments

- $H = \text{diag}(e^{-u})A \text{diag}(e^{-v})$  : Hungarian scaling.
- Row and column scalings with DGEEQU,  $B = D_R A D_C$ .
- $\min_{D_1, D_2 \in \mathcal{D}_n} \kappa_2(D_1 A D_2) \leq \rho(|A| |A^{-1}|) \leq n \min_{D_1, D_2 \in \mathcal{D}_n} \kappa_2(D_1 A D_2)$ ,  
[Rump' 03].

Problem	$\kappa_2(A)$	$\kappa(B)$	$\kappa_2(H)$	$\rho( A   A^{-1} )$
d_dyn	7.4e+6	8.4e+2	1.6e+1	6.2e+0
rotor1	2.4e+12	1.6e+8	3.0e+1	1.3e+0
lms_131	1.3e+15	1.3e+6	6.0e+1	1.8e+1
west0132	4.2e+11	2.8e+6	2.1e+2	6.5e+1
nnc261	2.9e+14	1.4e+12	4.5e+4	1.4e+4

## Theorem (Hook'14)

Let  $u, v \in \mathbb{R}^n$  and  $M \in \mathbb{R}_{\max}^{n \times n}$ . The matrix

$$\text{diag}(-u) \otimes M \otimes \text{diag}(-v)$$

has all of its max-plus singular values equal to zero iff  $(u, v)$  is a Hungarian pair for  $M$ .

Now if  $A \in \mathbb{C}^{n \times n}$  has entries which varies a lot in magnitude and  $(u, v)$  is a Hungarian pair for  $M = \mathcal{V}[A] \in \mathbb{R}_{\max}^{n \times n}$  then (heuristically)

$$\text{diag}(e^{-u})A \text{diag}(e^{-v})$$

should have its singular values map to  $\exp(0) = 1$ .

# Summary

- ▶ Max-plus roots, eigenvalues and singular values can offer order of magnitude approximations to their “classical” analogs for problems with large variations in the magnitude of the data.
- ▶ Plan to use tropical algebra to produce a polynomial eigensolver with better numerical stability property than eigensolvers such as `polyeig`.
- ▶ Investigate the effect of Hungarian scaling on eigenvalue condition numbers and backward errors.