

Solutions 2

3.10 Exercise

Prove that any embedding can be decomposed into an isomorphism and an identity embedding. I.e. show that if $\mathcal{A}, \mathcal{B} \in \mathcal{K}_\sigma$ and $\pi: \mathcal{A} \hookrightarrow \mathcal{B}$, then there exists $\mathcal{A}^* \in \mathcal{K}_\sigma$ such that $\mathcal{A}^* \subseteq \mathcal{B}$ and $\pi: \mathcal{A} \cong \mathcal{A}^*$.

Solution

Let (with the usual notation for \mathcal{A}, \mathcal{B}) $A^* := \{\pi(a) : a \in A\}$.

Then A^* is closed under each g_j ($j \in J$). For if

$\pi(a_1), \dots, \pi(a_{n(j)}) \in A^*$ (for $a_1, \dots, a_{n(j)} \in A$) then
 $g_j(\pi(a_1), \dots, \pi(a_{n(j)})) = \pi(f_j(a_1, \dots, a_{n(j)}))$ (by 2.4(c) since

$\pi: \mathcal{A} \hookrightarrow \mathcal{B}$). Since $f_j(a_1, \dots, a_{n(j)}) \in A$ this shows that
 $g_j(\pi(a_1), \dots, \pi(a_{n(j)})) \in A^*$.

A similar argument shows that $d_k \in A^*$ for each $k \in K$. Thus, A^* is the domain of a unique substructure, call it \mathcal{A}^* , of \mathcal{B} by 3.3. The relations, functions and distinguished elements of \mathcal{A}^* are necessarily given by 3.1 (with A^* in place of A) from which it immediately follows that $\pi: \mathcal{A} \hookrightarrow \mathcal{A}^*$ (since $\pi: \mathcal{A} \hookrightarrow \mathcal{B}$). But π obviously maps onto A^* . So $\pi: \mathcal{A} \cong \mathcal{A}^*$ as required.

4.5 Lemma

With the usual notation, if $\bar{a}, \bar{b} \in A^{\omega}$ (say $\bar{a} = \langle a_1, a_2, \dots \rangle$, $\bar{b} = \langle b_1, b_2, \dots \rangle$) and $a_p = b_p$ for all those p such that v_p occurs in τ , then
 $\tau^{\mathcal{A}}[\bar{a}] = \tau^{\mathcal{A}}[\bar{b}]$.

Proof.

We use induction on τ . If τ is c_k then $\tau^a[\bar{a}] = c_k = \tau^a[\bar{b}]$. If τ is v_p then $\tau^a[\bar{a}] = a_p$ and $\tau^a[\bar{b}] = b_p$. But v_p occurs in v_p so by the lemma hypothesis we have $a_p = b_p$, i.e. $\tau^a[\bar{a}] = \tau^a[\bar{b}]$.

Now suppose that τ is $f_j(\tau_1, \dots, \tau_{\mu(j)})$ (for some $j \in J$) and that the lemma is true for $\tau_1, \dots, \tau_{\mu(j)}$.

Suppose that $a_p = b_p$ for each p such that v_p occurs in τ .

Let $1 \leq r \leq \mu(j)$. Now for any p , if v_p occurs in τ_r then it occurs in τ (by definition of "occurring" — see bottom of page 14), so $a_p = b_p$. Hence, by the inductive hypothesis, $\tau_r^a[\bar{a}] = \tau_r^a[\bar{b}]$.

This is true for each $r = 1, \dots, \mu(j)$.

Therefore, $f_j(\tau_1^a[\bar{a}], \dots, \tau_{\mu(j)}^a[\bar{a}]) = f_j(\tau_1^a[\bar{b}], \dots, \tau_{\mu(j)}^a[\bar{b}])$.

I.e. $\tau^a[\bar{a}] = \tau^a[\bar{b}]$ (by 4.4 (iii)), as required.

4.8 Exercise

Describe in a mathematical way the collection of functions $\{\tau^a : A \rightarrow A : \tau \in J, \}$, where J denotes the set of terms in which at most the variable v_i occurs, in the following cases:

(1) $\mathcal{A} = \langle \mathbb{R}; +, \cdot, -; 0, 1 \rangle$.

(2) $\mathcal{A} = \langle \mathbb{Z}; \cdot, f \rangle$ where $f(z) = 2z$ (for $z \in \mathbb{Z}$).

Solution

(1) Here, $\tau^a[x]$ has the form $a_n x^n + a_{n-1} x^{n-1} + \dots + a_0$ for some $n \geq 0$ and $a_0, \dots, a_n \in \mathbb{R}$. I.e. $\tau^a[x]$ is a

polynomial in x with integer coefficients.

For let F_1, F_2 be the binary function symbols of \mathcal{L}_0 (interpreted by $+$, \cdot respectively) and F_3 the unary function symbol (interpreted by $-$). Let c_0, c_1 be the constant symbols (interpreted by $0, 1$) \ast we use

To see that $\tau^n[x]$ has the form \ast we use induction on τ .

We first consider the case that τ is v_1, c_0 or c_1 . We have $v_1^n[x] = x$, $c_0^n[x] = 0$ and $c_1^n[x] = 1$ which are all of the form \ast (with $n=1, a_1=1, a_0=0$; $n=0, a_0=0$; $n=0, a_0=1$ respectively).

Also if $\tau_1^n[x], \tau_2^n[x]$ are of the form \ast then $F_1(\tau_1, \tau_2)^n[x] = \tau_1^n[x] + \tau_2^n[x]$ can also be put into this form. Similarly for the inductive cases corresponding to F_2 and F_3 .

For the converse, for $r \geq 0$, let λ_r denote the term defined inductively by

$$\begin{aligned} \lambda_0 &:= c_1 \\ \lambda_{r+1} &= F_2(v_1, \lambda_r). \end{aligned}$$

Then $\lambda_r^n[x] = x^r$ (for all $x \in \mathbb{R}$).

Also define, $\delta_1 := c_1$, $\delta_{r+1} = F_1(c_1, \delta_r)$.

Then $\delta_r^n[x] = r$.

Hence $F_2(\delta_r, \lambda_s)^n[x] = r \cdot x^s$, and

$$F_3(F_2(\delta_r, \lambda_s))^n[x] = -r \cdot x^s.$$

Now we obtain \ast easily by induction on n ,

by use of F_1 .

(2) Here, a similar inductive argument shows that the only functions that we can get are of the form $x \mapsto 2^r x^s$ for some $r, s \in \mathbb{Z}$ with $r \geq 0$ and $s > 0$.

Further, all such functions are obtainable by repeated application of F_2 (multiplication) and F_1 (for f).
